









# Hidden Markov models

- ➢ Applications: Sequence alignment, gene detection et cetera
- ➢ Originated in speech recognition

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

# **CG-islands**

- ➢ Given 4 nucleotides: probability of occurrence is ~ 1/4. Thus, probability of occurrence of a dinucleotide is ~ 1/16.
- ➢ However, the frequencies of dinucleotides in DNA sequences vary widely.
- ➤ In particular, CG is typically underepresented (frequency of CG is typically < 1/16)</p>



- CG is the least frequent dinucleotide because C in CG is easily *methylated* and has the tendency to mutate into T afterwards
- However, the methylation is suppressed around genes in a genome. So, CG appears at relatively high frequency within these CG islands
- So, finding the *CG* islands in a genome is an important problem (gene finding)

#### CG-islands and the "Fair Bet Casino"

- ➢ The CG islands problem can be modeled after a problem named "The Fair Bet Casino"
- The game is to flip coins, which results in only two possible outcomes: Head or Tail
- The Fair coin will produce Heads and Tails with the same probability 1/2
- The **B**iased coin will produce **H**eads with prob.  $\frac{3}{4}$

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

# The "Fair Bet Casino" (cont'd)

Thus, we define the probabilities:

 $-P(H|F) = P(T|F) = \frac{1}{2}$ 

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics I

- $-P(H | B) = \frac{3}{4}, P(T | B) = \frac{1}{4}$
- -The crooked dealer changes between Fair and Biased coins with probability 0.1

# The Fair Bet Casino problem

- > Input: A sequence  $\mathbf{x} = x_1 x_2 x_3 \dots x_n$  of coin tosses (either *H* or *T*) made by two possible coins (*F* or *B*)
- ▶ Output: A sequence  $\pi = \pi_1 \pi_2 \pi_3 \dots \pi_m$ , with each  $\pi_i$  being either *F* or *B* indicating that  $x_i$  is the result of tossing the Fair or Biased coin, respectively
- Problem: Any observed outcome of coin tosses could have been generated by any sequence of states!



# P(x|fair coin) vs. P(x|biased coin)

- $P(\mathbf{x} | \text{ fair coin}) = P(x_1 \dots x_n | \text{ fair coin}) = \Pi_{i=1,n} p(x_i | \text{ fair coin}) = (1/2)^n$
- $P(\mathbf{x}| \text{ biased coin}) = P(x_1...x_n| \text{ biased coin}) = \Pi_{i=1,n} p(x_i| \text{ biased coin}) = (3/4)^k (1/4)^{n-k} = 3^k/4^n$ where k is the number of **H**eads in **x** (and *n-k* is the number of **T**ails)

















# HMM parameters

 $\Sigma$ : set of emission characters

 $\Sigma = \{H, T\}$  for coin tossing  $\Sigma = \{A, C, G, T\}$  for the CG-island problem

Q: set of hidden states, each emitting symbols from  $\Sigma$ Q={F,B} for coin tossing Q={CG-island, not CG-island} for the CGisland problem

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

# HMM Parameters (cont'd)

 $A = (a_k): a |Q| x |Q| matrix of probability of changing from state k to state l - transition probabilities$ 

 $E = (e_k(b)): a |Q| x |\Sigma| matrix of probability of$ emitting symbol b while being in state k- emission probabilities







# $P(x,\pi)$ Calculation

 $P(x,\pi) = P(x | \pi) P(\pi)$ : Probability that sequence *x* was generated by the path  $\pi$ :

$$P(\mathbf{x}, \boldsymbol{\pi}) = P(\pi_0 \rightarrow \pi_1) \cdot \prod_{i=1}^n P(x_i \mid \pi_i) \cdot P(\pi_i \rightarrow \pi_{i+1})$$
$$= a_{\pi_0, \pi_1} \cdot \prod_{i=1}^n e_{\pi_i} (x_i) \cdot a_{\pi_b, \pi_{i+1}}$$
$$= \prod_{i=0}^n e_{\pi_{i+1}} (x_{i+1}) \cdot a_{\pi_b, \pi_{i+1}}$$

where  $\pi_0$  and  $\pi_{n+1}$  are fictitious initial and terminal states *begin* and *end* TR Hundstern (MR304: Discrete structures for bioinformatics II 24



to solve the *Decoding Problem* > Every choice of  $\pi = \pi_1 \dots \pi_n$  corresponds to a

path in the graph

> The only valid direction in the graph is *eastward* 

Building an edit graph for the

**Decoding problem** 

Andrew Viterbi used the Manhattan grid model

This graph has  $|Q|^2(n-1)$  edges

















#### Decoding problem and Dynamic programming

Define  $s_{k,i}$  as the probability of emitting the prefix  $x_1 \dots x_i$  and reaching the state k

 $s_{l,i+1} = \max_{k \in \mathcal{Q}} \{s_{k,i} : \text{weight of edge between } (k,i) \text{ and } (l,i+1) \} =$ 

 $\max_{k \in \mathcal{Q}} \{ s_{k,i} \cdot a_{kl} \cdot e_l(x_{i+1}) \} =$ 

$$e_l(x_{i+1}) \cdot \max_{k \in \mathcal{O}} \{s_{k,i} \cdot a_{kl}\}$$

# Decoding problem (cont'd)

- ► Initialization:
- $-s_{begin,0} = 1$ -  $s_{k,0} = 0$  for  $k \neq begin$ .

Let  $\pi^*$  be the optimal path. Then,

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

$$P(\mathbf{x}, \boldsymbol{\pi}^*) = \max_{k \in O} \{s_{k,n} \cdot a_{k,end}\}$$

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

39



# Forward-backward problem

**Given:** a sequence of coin tosses generated by an HMM

**Goal:** find the probability that the dealer was using a biased coin at a particular time

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

# Forward algorithm

Define f<sub>k,i</sub> (forward probability) as the probability of emitting the prefix x<sub>1</sub>...x<sub>i</sub> and reaching the state π = k
 The recurrence for the forward algorithm:

/13

$$f_{k,i} = e_k(x_i) \cdot \sum f_{l,i-1} \cdot a_{lk}$$

Remember the Viterbi algorithm:  

$$s_{l,i+1} = e_l(x_{i+1}) \cdot \max_{k \in O} \{s_{k,i} \cdot a_{k}\}$$





# Backward-forward algorithm

The probability that the dealer used a biased coin at any moment *i*:

$$P(\pi_{i} = k | x) = \frac{P(x, \pi_{i} = k)}{P(x)} = \frac{f_{ki} \cdot b_{ki}}{P(x)}$$

where:

 $P(\mathbf{x}, \pi_i = k) = \sum_{\pi \text{ with } \pi_i = k} P(\mathbf{x}, \pi)$ : the sum of probabilities of all paths with  $\pi_i = k$  $P(\mathbf{x}) = \sum_{\pi} P(\mathbf{x}, \pi)$ : the sum of probabilities over all paths





- ► Probability of sequence and path: -  $P(\mathbf{x}, \pi) = P(\mathbf{x} | \pi) P(\pi) = \prod_{i} e_{\pi_{i+1}} (x_{i+1}) \cdot a_{\pi_{i}, \pi_{i+1}}$
- Viterbi gives the path  $\pi^*$  that maximizes  $P(\mathbf{x}, \pi)$ : -  $s_{l,i+1} = e_l(\mathbf{x}_{i+1}) \cdot \max_{k \in \mathbb{Q}} \{s_{k,i} \cdot a_{k,i}\}$
- ► Forward algorithm sums the probability of  $\mathbf{x}$  over all paths:  $- \sum_{\alpha} P(\mathbf{x}, \vec{\sigma}) = \sum_{\alpha} P(\mathbf{x} | \pi) P(\pi) = P(\mathbf{x})$   $- f_{k,l} = e_k(x_l) \cdot \sum_{l} f_{l,l}, a_{l,k}$ 
  - J.K,I "R("V"-1,J,I,I-I"")R
- Forward+backwards sums the probability of *x* over all paths with π<sub>i</sub> = k - Σ<sub>xwinh π<sub>i</sub> = k</sub> P(x,π)

40

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II





#### Finding distant members of a protein family

- A distant cousin of functionally related sequences in a protein family may have weak pairwise similarities with each member of the family and thus fail significance test using e.g. BLAST
- However, they may have weak similarities with *many* members of the family
- The goal is to align a sequence to all members of the family at once
- Family of related proteins can be represented by their multiple alignment and the corresponding profile



Aligned DNA sequences can be represented by a *4*·*n* profile matrix reflecting the frequencies of nucleotides in every aligned position.

Protein family can be represented by a  $20 \cdot n$  profile representing frequencies of amino acids.

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

53

# **HMMs**

- HMMs can also be used for aligning a sequence against a protein family
- Conserved positions in the family corresponds to *n* sequentially linked *match* states  $M_p,...,M_n$  in the profile HMM
- HMMs handle gaps better than profiles do

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

# Building a profile HMM

- > Multiple alignment is used to construct the HMM model
- Assign each column to a Match state in HMM. Add Insertion and Deletion state
- Estimate the emission probabilities according to amino acid counts in columns
- Estimate the transition probabilities between *Match, Deletion* and *Insertion* states







# Emission probabilities for insertions

Probability of emitting a symbol a at an insertion state  $I_i$ :

$$e_{Ii}(a) = p(a)$$

where p(a) is the frequency of the occurrence of the symbol *a* in all the sequences.

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics I

# Profile HMM alignment

➤ Define  $v_{j}^{M}(i)$  as the logarithmic likelihood score of the best path for matching  $x_{j}...x_{i}$  to a profile HMM ending with  $x_{i}$  emitted by the state  $M_{i}$ 

 $\succ v_{i}^{I}(i)$  and  $v_{i}^{D}(i)$  are defined similarly





# Making a collection of HMM for protein families

- Use Blast to separate a protein database into families of related proteins
- Construct a multiple alignment for each protein family
- Construct a profile HMM model and optimize the parameters of the model (transition and emission probabilities)
- Align the target sequence against each HMM to find the best fit between a target sequence and an HMM

62

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

#### Pfam

- Pfam decribes protein domains
- Each protein domain family in Pfam has:
  - Seed alignment: manually verified multiple alignment of a representative set of sequences
  - HMM built from the seed alignment for further database searches
  - Full alignment generated automatically from the HMM
- The distinction between seed and full alignments facilitates Pfam updates
  - Seed alignments are stable resources
  - Full alignments can be updated with newly found amino acid sequences

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

# Pfam

Pfam uses a tool called HMMER with the following architecture:





# HMM parameter estimation

- So far, we have (mostly) assumed that the transition and emission probabilities are known
- However, in most HMM applications, the probabilities are not known. It's very hard to estimate the probabilities.

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

#### HMM parameter estimation problem

#### Given

HMM with states and alphabet (emission characters)

lndependent training sequences  $x^1, \ldots x^m$ 

Find HMM parameters  $\Theta$  (that is,  $a_{kb} e_k(b)$ ) that maximize

# $P(x^1, \ldots, x^m \mid \Theta),$

the joint probability of the training sequences.

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

# Maximize the likelihood

 $P(x^{1}, ..., x^{m} | \boldsymbol{\Theta})$  as a function of  $\boldsymbol{\Theta}$  is called the likelihood of the model

The training sequences are assumed independent, therefore

 $P(x^{1}, \ldots, x^{m} \mid \boldsymbol{\Theta}) = \boldsymbol{\Pi}_{i} P(x^{i} \mid \boldsymbol{\Theta})$ 

The parameter estimation problem seeks  $\Theta$  that realizes

$$\max_{\Theta} \quad \prod_{i} P(x^{i} \mid \Theta)$$

In practice the log likelihood is computed to avoid underflow errors



Known paths for training sequences

- > CG islands are marked on training sequences
- One evening the casino dealer allows us to see when he changes dice
- ➤ A multiple alignment of the protein family is given

#### Unknown paths

- CG islands are not marked
- Do not see when the casino dealer changes dice
- A multiple alignment of the protein family is *not* given

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

#### Known paths

 $A_{kl}$  = # of transitions from state k to state l in the training sequences

 $E_k(b) = \#$  of times *b* is emitted from state *k* in the training sequences

Compute  $a_{kl}$  and  $e_k(b)$  as maximum likelihood estimators:

$$a_{kl} = A_{kl} / \sum_{l' \in Q} A_{kl'}$$
$$e_k(b) = E_k(b) / \sum_{b' \in \Sigma} E_k(b')$$

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

# **Pseudocounts**

- Some state *k* may not appear in any of the training sequences. This means  $A_{kl} = 0$  for every state *l* and  $a_{kl}$  cannot be computed with the given equation
- To avoid this overfitting use predetermined pseudocounts r<sub>kl</sub> and r<sub>k</sub>(b).

 $A_{kl} = \#$  of transitions  $k \rightarrow l + r_{kl}$ 

 $E_k(b) = \#$  of emissions of b from  $k + r_k(b)$ 

The pseudocounts reflect our prior biases about the probability values

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

# Unknown paths: Viterbi training

<u>Idea</u>: use Viterbi decoding to compute the most probable paths for training sequences x

- Start with some guess for initial parameters
- Iterate :
  - Compute the most probable path π\* for each x using the current parameters (Viterbi algorithm)
     Stop if no change in π\*
  - Determine  $A_{kl}$  and  $E_k(b)$  using the computed paths
  - Compute new parameters a<sub>kl</sub> and e<sub>k</sub>(b) using the same formulas as before



73

# Unknown paths: Baum-Welch

#### Idea:

- 1. Guess initial values for parameters. art and experience, not science
- 2. Estimate new (better) values for parameters.
- 3. Repeat until stopping criteria is met. what criteria ?

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

#### Better values for parameters

Would need the  $A_{kl}$  and  $E_{k}(b)$  values but cannot count (the path is unknown) and do not want to use a most probable path

For all states k,l, symbol b and training sequence x

Compute  $A_{kl}$  and  $E_k(b)$  as expected values, given the current parameters

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

#### Notation

For any sequence of characters x emitted along some <u>unknown path</u>  $\pi$ , denote by  $\pi_i = k$  the assumption that the state at position i(in which  $x_i$  is emitted) is k

# **Probabilistic setting for** $A_{k,l}$

Given  $x^1, \ldots, x^m$  consider a discrete probability space with elementary events

 $\varepsilon_{k,l} = k \rightarrow l$  is taken in  $x^1, \dots, x^m$ "

For each x in  $\{x^i, ..., x^m\}$  and each position i in x let  $Y_{x,i}$  be a random variable defined by

 $Y_{x,i}(\varepsilon_{k,l}) = \begin{cases} 1, & \text{if } \pi_i = k \text{ and } \pi_{i+1} = l \\ 0, & \text{otherwise} \end{cases}$ Define  $Y = \sum_x \sum_i Y_{x,i}$  random variable that counts # of times the event  $\varepsilon_{k,l}$  happens in  $x^l, \dots, x^m$ . I.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II 77

# The meaning of $A_{kl}$

Let  $A_{kl}$  be the expectation of Y

 $E(Y) = \sum_{x} \sum_{i} E(Y_{x,i}) = \sum_{x} \sum_{i} P(Y_{x,i} = 1) =$  $\sum_{x} \sum_{i} P(\{\varepsilon_{k,l} \mid \pi_{i} = k \text{ and } \pi_{i+1} = l\}) =$  $\sum_{x} \sum_{i} P(\pi_{i} = k, \pi_{i+1} = l \mid x)$ 

Need to compute  $P(\pi_i = k, \pi_{i+1} = l \mid x)$ 

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

# Probabilistic setting for $E_k(b)$

Given  $x^1, \ldots, x^m$  consider a discrete probability space with elementary events

 $\varepsilon_{k,b} = "b$  is emitted in state k in  $x^{i}, \dots, x^{m}$ " For each x in  $\{x^{i}, \dots, x^{m}\}$  and each position i in x let  $Y_{x,i}$  be a random variable defined by

$$Y_{x,i}(\varepsilon_{k,b}) = \begin{cases} 1, & \text{if } x_i = b \text{ and } \pi_i = k \\ 0, & \text{otherwise} \end{cases}$$

Define  $Y = \sum_{x} \sum_{i} Y_{x,i}$  random variable that counts # of times the event  $\varepsilon_{k,b}$  happens in  $x^{1}, \dots, x^{m}$ . The meaning of  $E_k(b)$ 

Let  $E_k(b)$  be the expectation of Y

$$E(Y) = \sum_{x} \sum_{i} E(Y_{x,i}) = \sum_{x} \sum_{i} P(Y_{x,i} = 1) =$$
  
$$\sum_{x} \sum_{i} P(\{\varepsilon_{k,b} \mid x_{i} = b \text{ and } \pi_{i} = k\}) =$$
  
$$\sum_{x} \sum_{i} P(\{\varepsilon_{k,b} \mid x_{i} = b, \pi_{i} = k\}) = \sum_{x} \sum_{i} P(\pi_{i} = k \mid x)$$

$$\sum_{x} \{i|x_i=b\} \quad ((a_{k,b}+x_i-b,x_i-b)) \quad \sum_{x} \{i|x_i=b\}$$

Need to compute  $P(\pi_i = k \mid x)$ 



# Compute $A_{kl}$ (1)

Prob  $k \rightarrow l$  is taken at position i of x $P(\pi_i = k, \pi_{i+1} = l \mid x_1 \dots x_n) = P(x, \pi_i = k, \pi_{i+1} = l) / P(x)$ 

Compute P(x) using either forward or backward values We'll show that  $P(x, \pi_i = k, \pi_{i+1} = l) = b_{li+1} \cdot e_l(x_{i+1}) \cdot a_{kl} \cdot f_{kl}$ 

Expected # times  $k \rightarrow l$  is used in training sequences  $A_{kl} = \sum_{x} \sum_{i} (b_{li+1} \cdot e_{i}(x_{i+1}) \cdot a_{kl} \cdot f_{ki}) / P(x)$ 

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

# Compute $A_{kl}$ (2)

$$\begin{split} P(\mathbf{x}, \pi_{i} = \mathbf{k}, \pi_{i+1} = \mathbf{l}) &= \\ P(\mathbf{x}_{1} \dots \mathbf{x}_{p}, \pi_{i} = \mathbf{k}, \pi_{i+1} = \mathbf{l}, \mathbf{x}_{i+1} \dots \mathbf{x}_{p}) &= \\ P(\pi_{i+1} = \mathbf{l}, \mathbf{x}_{i+1} \dots \mathbf{x}_{n} \mid \mathbf{x}_{1} \dots \mathbf{x}_{p}, \pi_{i} = \mathbf{k}) \cdot P(\mathbf{x}_{1} \dots \mathbf{x}_{p}, \pi_{i} = \mathbf{k}) = \\ P(\pi_{i+1} = \mathbf{l}, \mathbf{x}_{i+1} \dots \mathbf{x}_{n} \mid \pi_{i} = \mathbf{k}) \cdot f_{ki} = \\ P(\mathbf{x}_{i+1} \dots \mathbf{x}_{n} \mid \pi_{i} = \mathbf{k}, \pi_{i+1} = \mathbf{l}) \cdot P(\pi_{i+1} = \mathbf{l} \mid \pi_{i} = \mathbf{k}) \cdot f_{ki} = \\ P(\mathbf{x}_{i+1} \dots \mathbf{x}_{n} \mid \pi_{i+1} = \mathbf{l}) \cdot \mathbf{a}_{kl} \cdot f_{ki} = \\ P(\mathbf{x}_{i+2} \dots \mathbf{x}_{n} \mid \mathbf{x}_{i+1}, \pi_{i+1} = \mathbf{l}) \cdot P(\mathbf{x}_{i+1} \mid \pi_{i+1} = \mathbf{l}) \cdot \mathbf{a}_{kl} \cdot f_{ki} = \\ P(\mathbf{x}_{i+2} \dots \mathbf{x}_{n} \mid \pi_{i+1} = \mathbf{l}) \cdot e_{\mathbf{l}}(\mathbf{x}_{i+1}) \cdot \mathbf{a}_{kl} \cdot f_{ki} = \\ P(\mathbf{x}_{i+2} \dots \mathbf{x}_{n} \mid \pi_{i+1} = \mathbf{l}) \cdot e_{\mathbf{l}}(\mathbf{x}_{i+1}) \cdot \mathbf{a}_{kl} \cdot f_{ki} = \\ D(\mathbf{x}_{i+1} \cdot \mathbf{e}_{\mathbf{l}}(\mathbf{x}_{i+1}) \cdot \mathbf{a}_{kl} \cdot f_{ki} = \\ D(\mathbf{x}_{i+1} \cdot \mathbf{e}_{\mathbf{l}}(\mathbf{x}_{i+1}) \cdot \mathbf{a}_{kl} \cdot f_{ki} = \\ D(\mathbf{x}_{i+1} \dots \mathbf{x}_{n} \mid \pi_{i+1} = \mathbf{l}) \cdot e_{\mathbf{l}}(\mathbf{x}_{i+1}) \cdot \mathbf{a}_{kl} \cdot f_{ki} = \\ D(\mathbf{x}_{i+1} \cdot \mathbf{e}_{\mathbf{l}}(\mathbf{x}_{i+1}) \cdot \mathbf{a}_{kl} \cdot f_{ki} = \\ D(\mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} = \mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} = \mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} = \\ D(\mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} = \mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} = \\ D(\mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} = \mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} + \\ D(\mathbf{x}_{i+1} \cdot \mathbf{x}_{i+1} - \mathbf{x}_{i+1} \cdot \mathbf{x}_{i+$$

Compute  $E_k(b)$ Prob  $x_i$  of x is emitted in state k  $P(\pi_i = k \mid x_1...x_n) = P(\pi_i = k, x_1...x_n)/P(x)$   $P(\pi_i = k, x_1...x_n) = P(x_1...x_p\pi_i = k, x_{i+1}...x_n) =$   $P(x_{i+1}...x_n \mid x_1...x_p\pi_i = k) \cdot P(x_1...x_p\pi_i = k) =$   $P(x_{i+1}...x_n \mid \pi_i = k) \cdot f_{ki} = b_{ki} \cdot f_{ki}$ Expected # times b is emitted in state k $E_k(b) = \sum_x \sum_{i:x_i=b} (f_{ki} \cdot b_{ki})/P(x)$ 





# The Baum-Welch algorithm

#### Initialization:

Pick the best-guess for model parameters (or arbitrary)

#### Iteration:

- 1. Forward for each x
- 2. Backward for each x
- 3. Calculate  $A_{kb} E_k(b)$
- 4. Calculate new  $a_{kb} e_k(b)$
- 5. Calculate new log-likelihood
- Until log-likelihood does not change much

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics II

# Baum-Welch analysis

Log-likelihood is increased by iterations Baum-Welch is a particular case of the EM (expectation maximization) algorithm

T.R. Hvidsten: 1MB304: Discrete structures for bioinformatics

Convergence to local maximum. Choice of initial parameters determines local maximum to which the algorithm converges

22